



MAFTIA WP2 Demonstration Transaction Service and Middleware

University of Newcastle upon Tyne

University of Lisboa

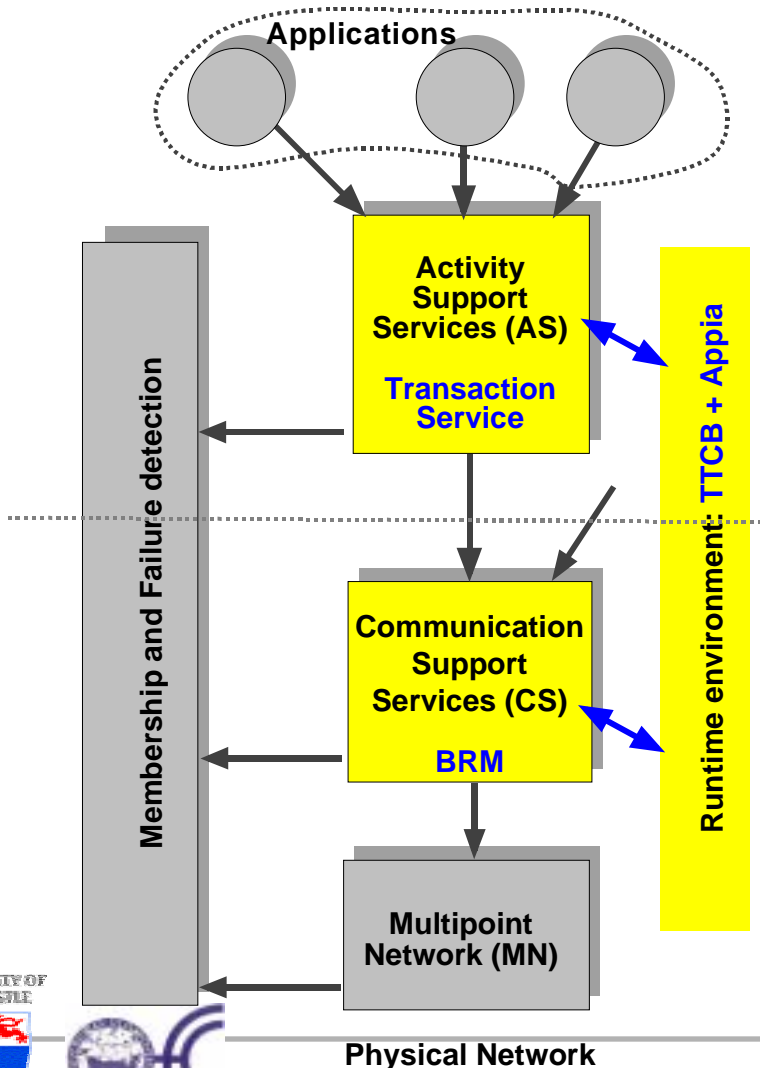
February 2003



Summary

- MAFTIA middleware
- TTCB and BRM
- Transaction Service
- Scenario: *an intrusion-tolerant transactional tic tac toe*
- Demonstration
 - Normal and multiparty transactions
 - Resilience to intrusions
 - Limits of intrusion tolerance

Maftia Middleware



- **Activity services**

- ☞ **Transaction Service**

- University of Newcastle upon Tyne (UK)

- **Communication services**

- ☞ Partially synchronous protocols (**BRM**)
University of Lisboa (P)

- ☞ Asynchronous protocols:
IBM, Zurich (CH)

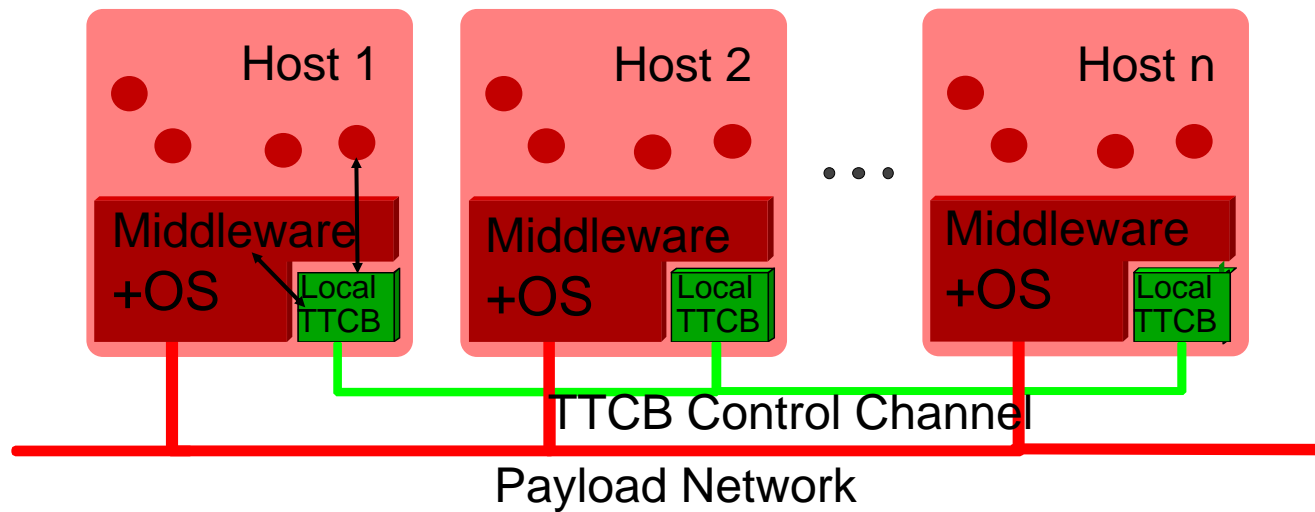
- **Runtime support**

- ☞ **TTCB and Appia**

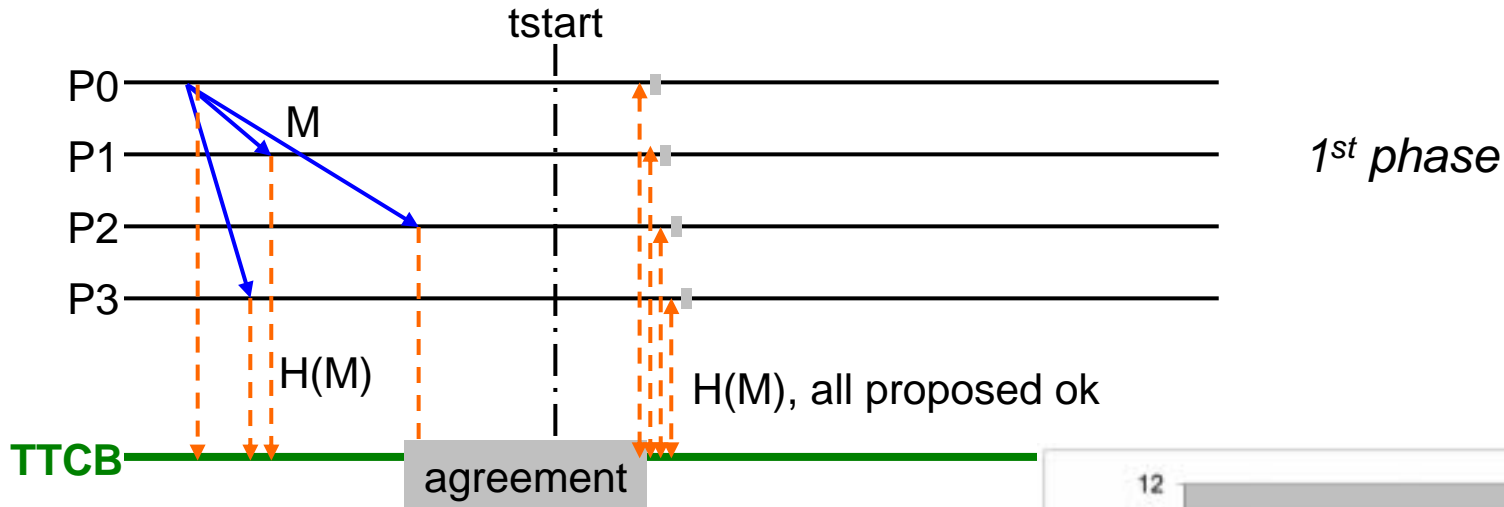
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TTCB Wormhole Model

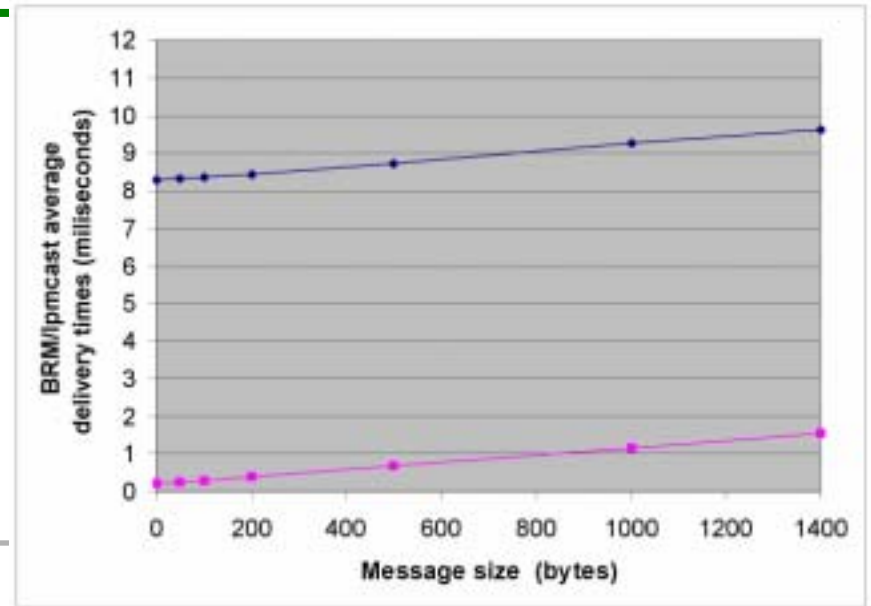
- Most of the system is Byzantine-on-failure except...
- TTCB wormhole: secure, real-time, limited services



Byzantine Reliable Multicast BRM



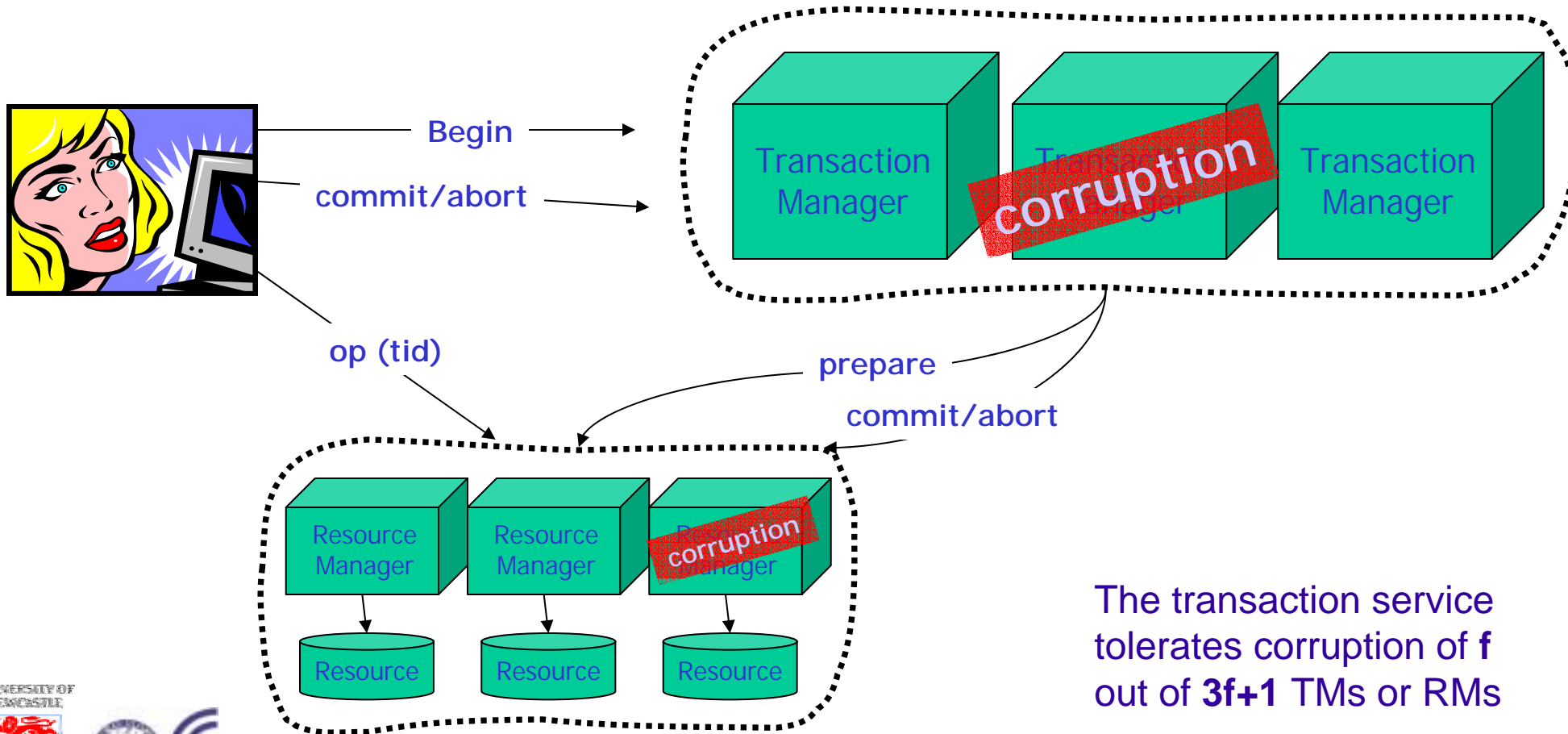
5-node delivery times



Transaction Service

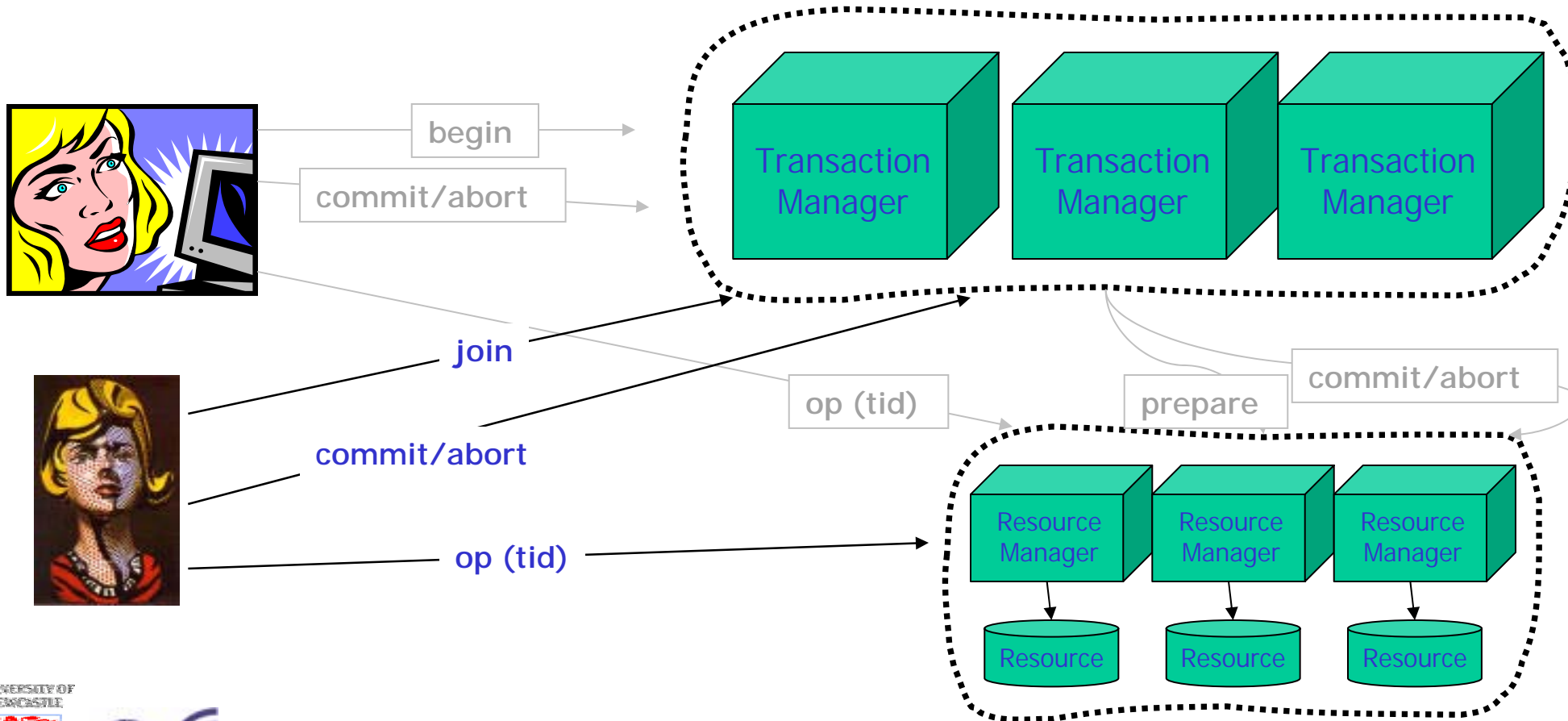
- A CORBA-style transaction service
- Provide ACID properties: atomicity, consistency, isolation, durability
- Multiparty transactions
- Intrusion tolerant

Intrusion Masking - Replication



The transaction service tolerates corruption of f out of $3f+1$ TMs or RMs

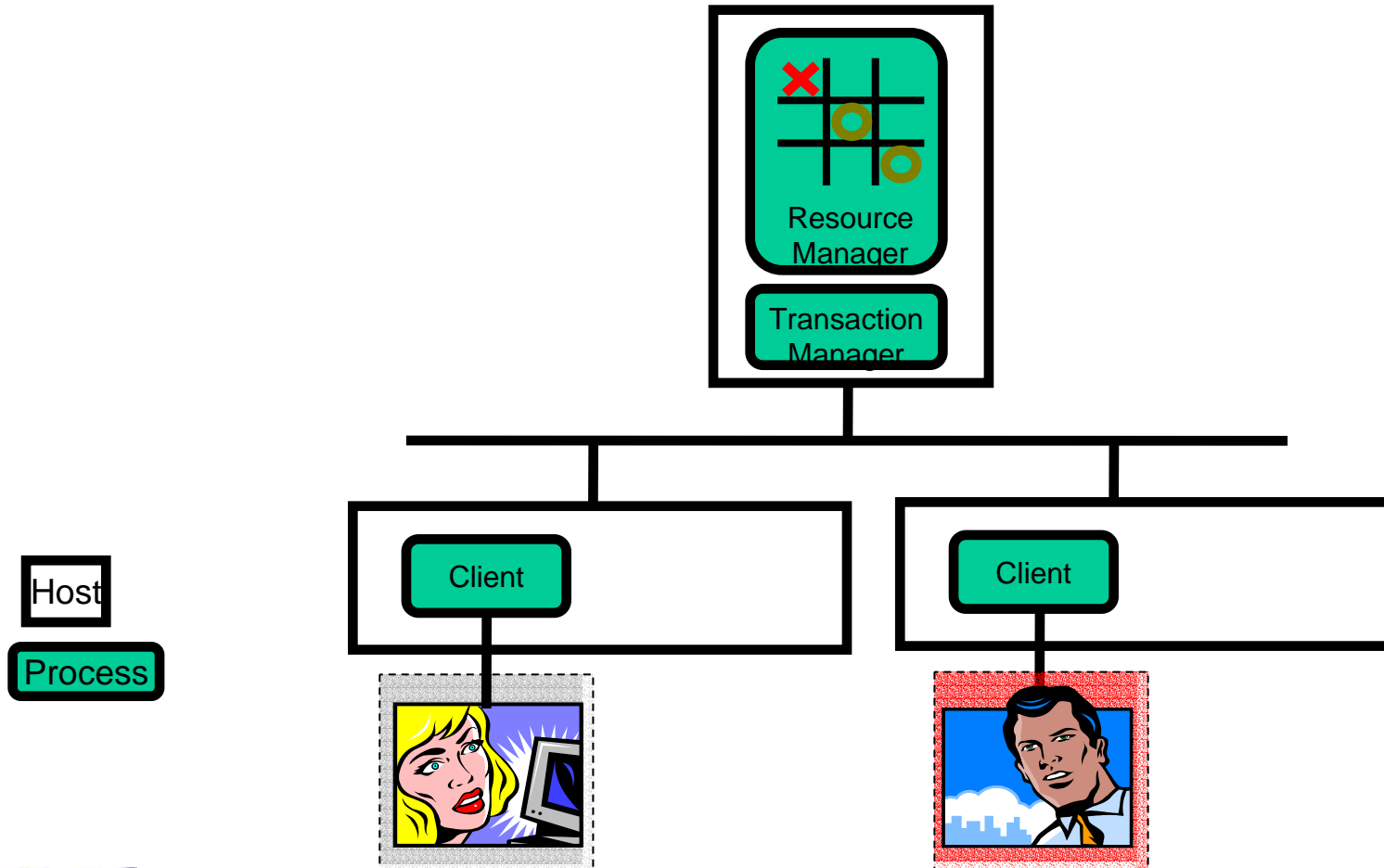
Multiparty Transactions



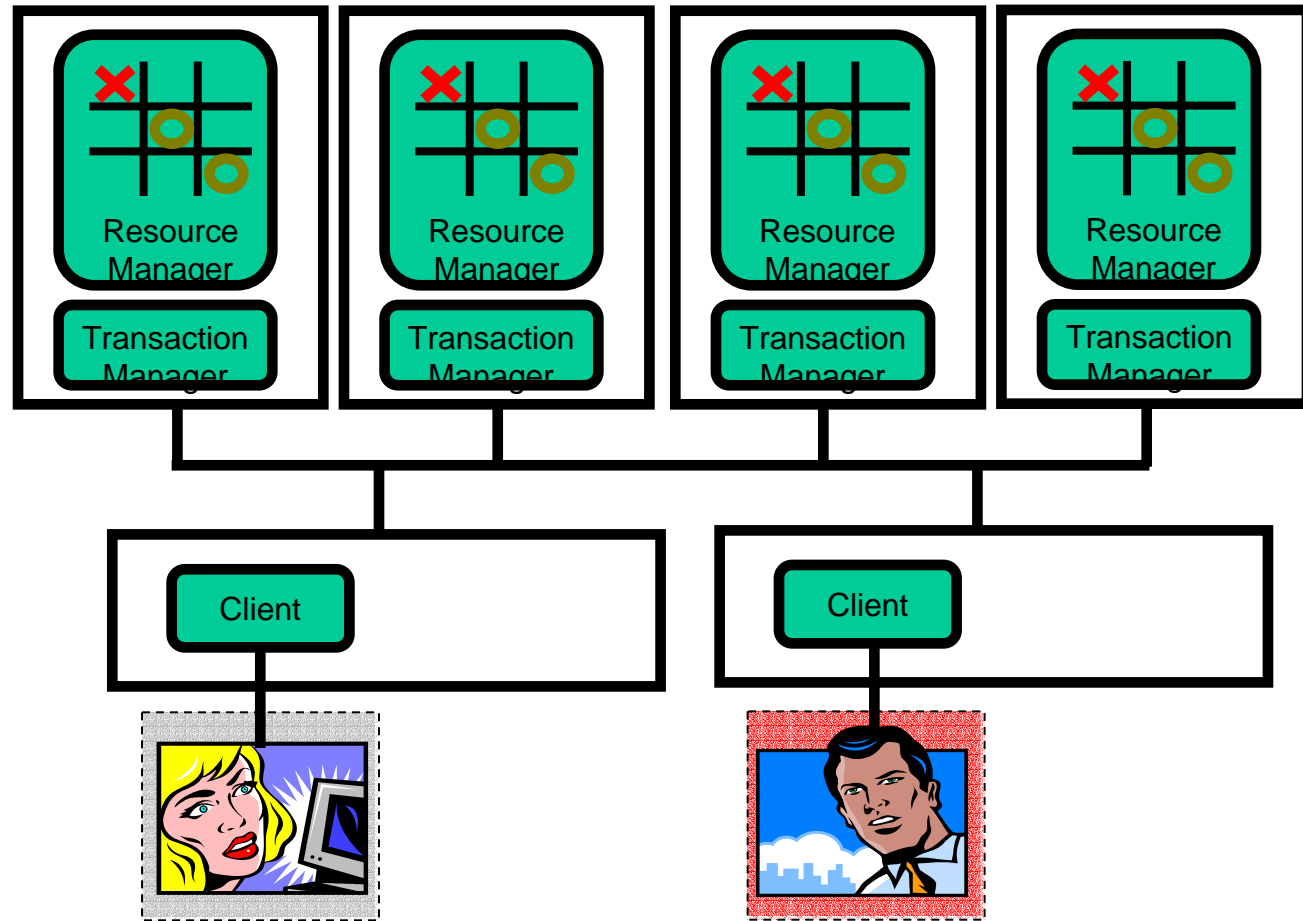
Demonstration Scenario

- A simple *tic tac toe* game
- There are two players who interact the game (transactional resource)
- We show that an intrusion tolerant transaction service can provide a correct service despite intrusions

Standard Transactional TicTacToe



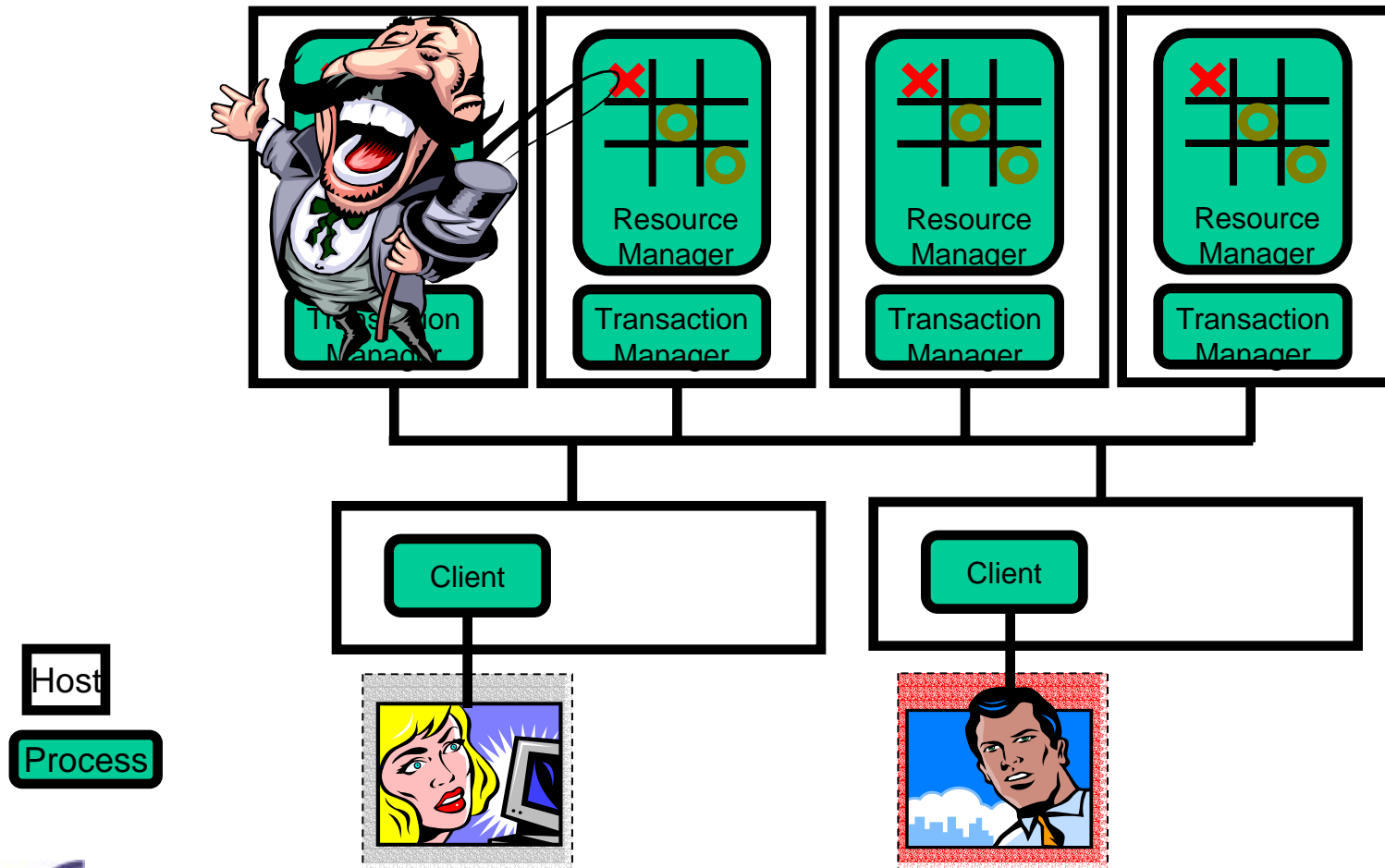
IT Transactional Tic Tac Toe



Hacker Behaviour

- A hacker may be able to install trapdoors in the TMs and RMs that allow him to bring them under his control
- Hacker wishes to influence the outcomes of games

Hacked Tic Tac Toe



Demonstration

- 1 resource manager (tic tac toe board) per machine (4 total)
- 1 transaction manager per machine (4 total)
- 2 clients
- TTCB: one local TTCB per machine and control network
- Appia and protocol stack (BRM)

Demonstration 1: Normal Case

- Making a move : Player begins a transaction, places a piece, and either aborts or commits the change making it visible to other players
- Benign Environment:
 - Player 1 makes move, player 2 makes a move
 - Player 1 makes move but aborts
 - Multiparty transaction: player 1 begins a transaction, player 2 joins, both make a move and commit

Demonstration 2: Intrusion Tolerance (RM)

- Hack one RM
 - Player 1 makes move
 - Hacked RM flips piece
 - Player 2 starts move and refreshes board
 - **The hacker is not able to change the moves**

Demonstration 3: Intrusion Tolerance (TM)

- Hack one TM
 - Player 1 makes move and commits
 - Hacked TM requests resource managers to abort
 - Move is committed despite intrusion
 - **The hacker is not able to change the decision**

Demonstration 4:

Limits of Intrusion Tolerance

- Hack two more RMs (3 total)
 - Hackers now control the majority of the RMs...
 - Player 1 makes move and commits
 - Three out of four RMs flip piece
 - Player 2 begins move and refreshes board state from resource managers
 - Board displays as intended by the hacker rather than Player 1
 - **To tolerate more intrusions we require more replicas (10 for 3 intrusions)**

Summary

- **Demonstrated IT transaction service**
 - Distributed CORBA style transaction service with support for multiparty transactions
 - Intrusion tolerance service composed of intrusion tolerant protocols
 - Can tolerate intrusions, the number is related to the degree of replication
- **Demonstrated**
 - Application of MAFTIA design principles (WP1) to Middleware activity service (WP2)
 - Integration between middleware layers and TTCB